

To adopt or not to adopt L4S-compatible congestion control?

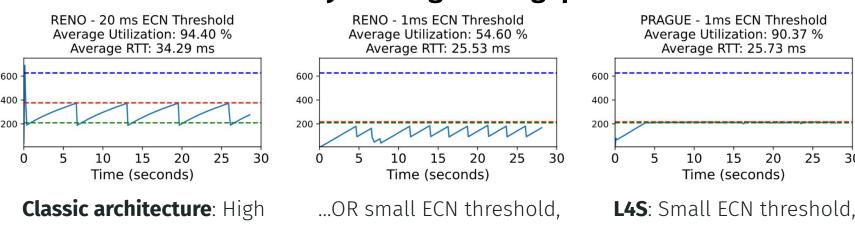


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What is L4S?

Low Latency, Low Loss, Scalable Throughput (L4S) [3] architecture is designed to enable both high throughput and low latency while coexisting with classic flows.

Classic congestion control cannot achieve extremely low delay with high throughput.



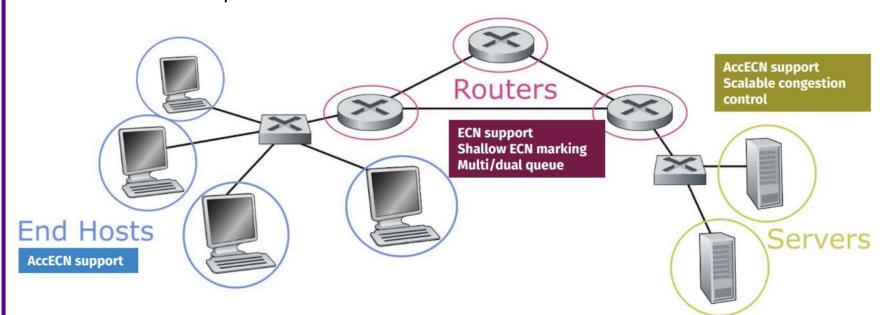
ECN threshold, high delay...

But: Unlike other mechanisms for low latency, L4S involves endpoints and middleboxes.

low throughput

low delay, high throughput

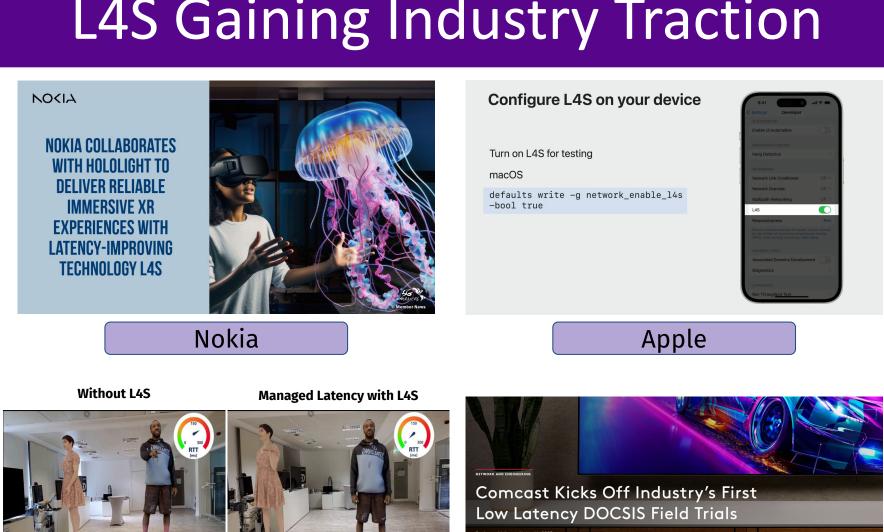
RQ-1



When all of these components are in place, an L4S flow can achieve high throughput with very low latency.

- → **However**, like any new Internet technology, the deployment of L4S will be **incremental**.
- → In the initial stages of deployment, L4S flows will coexist with classic flows at L4S or non-L4S bottleneck routers.

L4S Gaining Industry Traction



Comcast

Ericsson

Research Questions

- Can senders be assured that **TCP Prague** (L4S flow) will not cause harm to, or be harmed by, another flow at a shared bottleneck link? [1], we partially analyzed in [2]
- Does L4S-compatible BBRv2 have more favorable properties for adoption than TCP Prague? [1]
- Is the harm caused by or to the L4S-compatible flow mitigated when the bottleneck is shared by a large number of flows? [1]
- 4. BBRv3 enables ECN based on a path delay threshold. Is this approach better suited for L4S-compatible congestion control deployment? [1]

Experiment Methodology

*with competing flows in parentheses

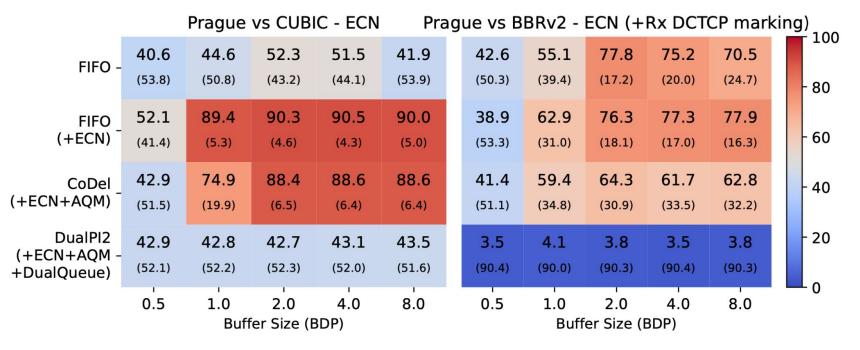
Line topology on **FABRIC** [4] testbed comprising two sending and two receiving hosts.

Residential home broadband scenario base delay by 10 milliseconds link capacity L4S Receiver L4S Sender **Delay Node** Classic Receiver Classic Sender

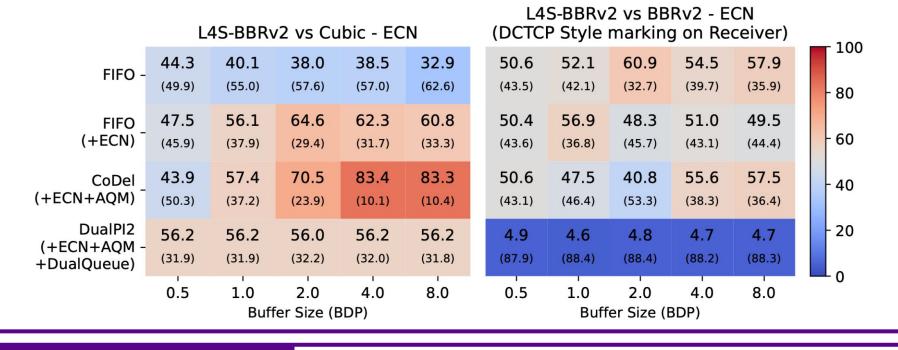
- TCP flow(s) from the L4S sender and from the classic sender, each for a duration of 60 seconds.
- A wide variety of queue types that may be encountered at the bottleneck router.
- FIFO, CoDel, PIE, FQ, L4S-aware FQ-CoDel, DualPI2
- TCP L4S flows: Prague and L4S-compatible BBRv2
- non-L4S flows: CUBIC, BBRv1/v2/v3

Findings - L4S Flows Throughput

→ Senders cannot be fully assured, as TCP Prague may cause harm or be harmed in various bottlenecks, though it coexists well in drop-based and fairness-enforcing queues.

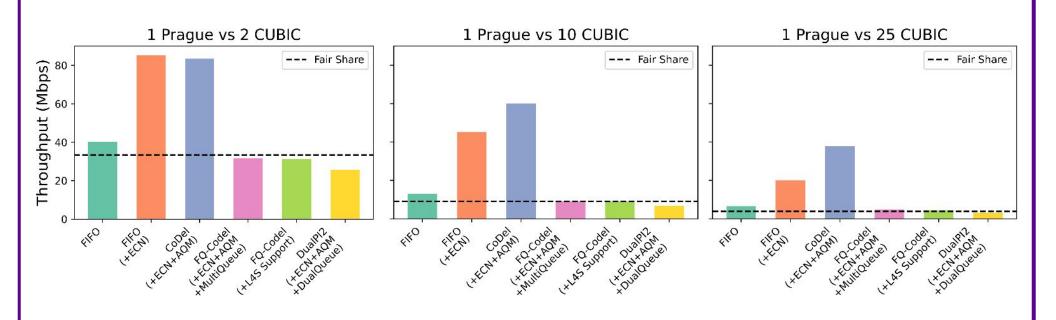


→ While L4S-compatible BBRv2 has more favorable properties for adoption than TCP Prague, senders still cannot be fully assured that it will not cause harm to, or be harmed by, other flows at a shared bottleneck link.

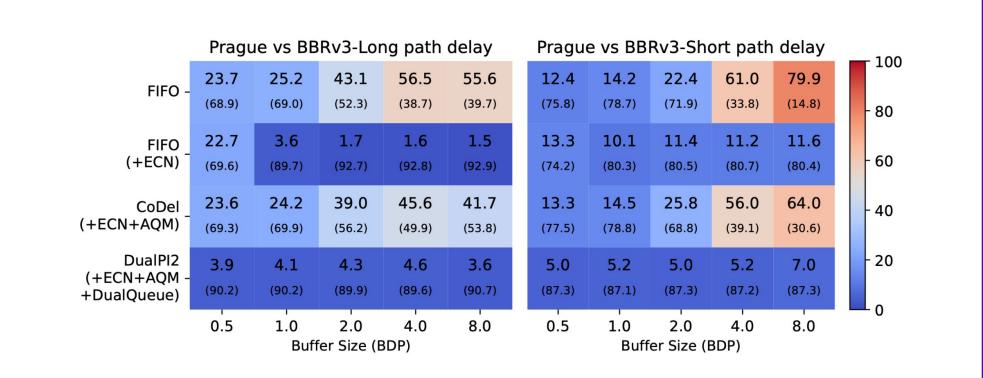


RQ-3

→ The harm is not necessarily mitigated when a scalable flow shares a bottleneck with multiple classic flows.



→ BBRv3 approach is not more favorable, as it presents coexistence challenges with both TCP Prague and L4S-compatible BBRv2 across various bottlenecks.



Takeaway

→ Given that the sender of an L4S flow cannot be sure what type of queue is at the bottleneck router or what other flows will share it, safe coexistence cannot be guaranteed right now.

Future Work

- →Investigating the impact on L4S adoption and strategies for mitigation more directly.
- →Extending the evaluation to more realistic traffic other and patterns network settings.
- [1] F. B. Sarpkaya, F. Fund, and S. Panwar, "To adopt or not to adopt L4S-compatible congestion control? Understanding performance in a partial L4S deployment," Under review, Oct. 2024.
- [2] F. B. Sarpkaya, A. Srivastava, F. Fund, and S. Panwar, "To switch or not to switch to TCP Prague? Incentives for adoption in a partial L4S deployment," ANRW '24, Vancouver, Canada, Jul. 2024. [QR]
- References [3] B. Briscoe, K. D. Schepper, M. Bagnulo, and G. White, "Low Latency, Low Loss, and Scalable Throughput (L4S) Internet Service: Architecture," RFC 9330, Jan. 2023. [Online].
 - [4] Baldin, Ilya, et al. "FABRIC: A national-scale programmable experimental network infrastructure." IEEE Internet Computing 23.6 (2019): 38-47.

